

# Asymptotics and Simulation of Heavy-Tailed Processes

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## Outline

## Contents

<b>1 Introduction to Rare-Event Simulation</b>	<b>1</b>
1.1 The Problem with Monte Carlo . . . . .	1
1.2 Importance Sampling . . . . .	2
<b>2 Importance Sampling in a Heavy-Tailed Setting</b>	<b>4</b>
2.1 The Conditional Mixture Algorithm for Random Walks . . . . .	4
2.2 Performance of the Conditional Mixture Algorithm . . . . .	5
<b>3 Markov Chain Monte Carlo in Rare-Event Simulation</b>	<b>6</b>
3.1 Efficient Sampling for a Heavy-Tailed Random Walk . . . . .	7
3.2 Efficient Sampling of Heavy-Tailed Random Sums . . . . .	8

## 1 Introduction to Rare-Event Simulation

### Simulation of Heavy-Tailed Processes

- Goal: improve on the computational efficiency of standard Monte Carlo.
- Design: The large deviations analysis lead to a heuristic way to design efficient algorithms.
- Efficiency Analysis: The large deviations analysis can be applied to theoretically quantify the computational performance of algorithms.

### 1.1 The Problem with Monte Carlo

#### The Performance of Monte Carlo

- Let  $X^1, \dots, X^N$  be independent copies of  $X$ .
- The Monte Carlo estimator is

$$\hat{p} = \frac{1}{N} \sum_{k=1}^N I\{X^k > b\}.$$

- Its standard deviation is

$$\text{std} = \frac{1}{\sqrt{N}} \sqrt{p(1-p)}$$

- To obtain a Relative Error (std/p) of 1% you need to take  $N$  such that

$$\frac{\frac{1}{\sqrt{N}} \sqrt{p(1-p)}}{p} \leq 0.01 \Leftrightarrow N \geq 10^4 \frac{1-p}{p}$$

## 1.2 Importance Sampling

### Introduction to Importance Sampling<sup>1</sup>

- The problem with Monte Carlo is that few samples hit the rare event.
- This problem can be fixed by sampling from a distribution which puts more mass on the rare event.
- Must compensate for not sampling from the original distribution.
- Key issue: How to select an appropriate sampling distribution?

**Monte Carlo and Importance Sampling The Basics for Computing**  $F(A) = \mathbb{P}\{X \in A\}$

#### Monte Carlo

- Sample  $X^1, \dots, X^N$  from  $F$ .
- Empirical measure

$$\mathbf{F}_N(\cdot) = \frac{1}{N} \sum_{k=1}^N \delta_{X^k}(\cdot).$$

- Plug-in estimator

$$\hat{p} = \mathbf{F}_N(A).$$

#### Importance Sampling

- Sample  $\tilde{X}^1, \dots, \tilde{X}^N$  from  $\tilde{F}$ .
- Weighted empirical measure

$$\tilde{\mathbf{F}}_N^w(\cdot) = \frac{1}{N} \sum_{k=1}^N \frac{dF}{d\tilde{F}}(\tilde{X}^k) \delta_{\tilde{X}^k}(\cdot).$$

- Plug-in estimator

$$\hat{p} = \tilde{\mathbf{F}}_N^w(A) = \frac{1}{N} \sum_{k=1}^N \frac{dF}{d\tilde{F}}(\tilde{X}^k) I\{\tilde{X}^k \in A\}.$$

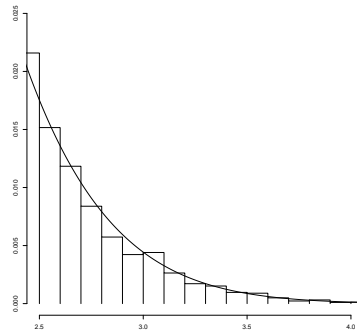
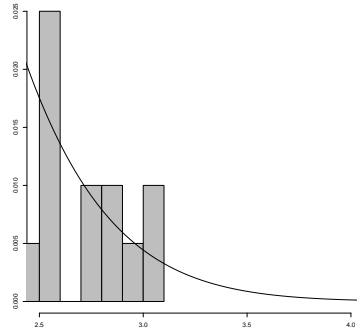
#### Importance Sampling

- The importance sampling estimator is unbiased:

$$\tilde{E} \left[ \frac{dF}{d\tilde{F}}(\tilde{X}) I\{\tilde{X} \in A\} \right] = \int_A \frac{dF}{d\tilde{F}} d\tilde{F} = \int_A dF = F(A).$$

- Its variance is

$$\begin{aligned} \text{Var} \left( \frac{dF}{d\tilde{F}}(\tilde{X}) I\{\tilde{X} \in A\} \right) &= \tilde{E} \left[ \left( \frac{dF}{d\tilde{F}} \right)^2 I\{\tilde{X} \in A\} \right] - F(A)^2 \\ &= \int_A \left( \frac{dF}{d\tilde{F}} \right)^2 d\tilde{F} - F(A)^2 \\ &= \int_A \frac{dF}{d\tilde{F}} dF - F(A)^2 \\ &= E \left[ \frac{dF}{d\tilde{F}} I\{X \in A\} \right] - F(A)^2. \end{aligned}$$



**Illustration Sampling the tail**  
**Monte Carlo Importance Sampling**

**Quantifying Efficiency Variance Based Efficiency Criteria**

- Basic idea: variance roughly the size of  $p^2$ .
- Embed in a sequence of problems such that  $p_n = \mathbb{P}\{X_n \in A\} \rightarrow 0$ .
- Logarithmic Efficiency: for some  $\epsilon > 0$

$$\limsup_{n \rightarrow \infty} \frac{\text{Var}(\hat{p}_n)}{p_n^{2-\epsilon}} < \infty.$$

- Strong Efficiency/Bounded Relative Error:

$$\sup_n \frac{\text{Var}(\hat{p}_n)}{p_n^2} < \infty.$$

- Vanishing relative error:

$$\limsup_n \frac{\text{Var}(\hat{p}_n)}{p_n^2} = 0.$$

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<sup>1</sup>See e.g. [1] for an introduction.

### The Zero-Variance Change of Measure

- There is a best choice of sampling distribution, the zero-variance change of measure, given by

$$F_A(\cdot) = \mathbb{P}\{X_n \in \cdot \mid X_n \in A\}.$$

- For this choice

$$\frac{dF}{dF_A}(x) = \mathbb{P}\{X_n \in A\}I\{X_n \in A\} = p_n I\{X_n \in A\},$$

and hence

$$\text{Var}(\hat{p}_n) = E[p_n I\{X_n \in A\}] - F(A)^2 = 0.$$

## 2 Importance Sampling in a Heavy-Tailed Setting

### Example Problem A Random Walk with Heavy Tails

- Let  $\{Z_k\}$  be iid non-negative and regularly varying (index  $\alpha$ ) with density  $f$ .
- Put  $S_m = Z_1 + \dots + Z_m$ .
- Compute  $\mathbb{P}\{S_m > n\}$ .
- The zero-variance sampling distribution is

$$\mathbb{P}\{(Z_1, \dots, Z_m) \in \cdot \mid S_m > n\}.$$

- By regular variation we know that

$$\mathbb{P}\{(Z_1, \dots, Z_m) \in \cdot \mid S_m > n\} \rightarrow \mu(\cdot),$$

where  $\mu$  is concentrated on the coordinate axis. But  $\mu$  and  $F$  are singular!

### 2.1 The Conditional Mixture Algorithm for Random Walks

#### The Conditional Mixture Algorithm<sup>2</sup>

- Suppose  $S_{i-1} = s$ . Sample  $Z_i$  as follows
  1. If  $s \geq n$ , sample  $Z_i$  from the original density  $f$ .
  2. If  $s \leq n$ , sample  $Z_i$  from the mixture

$$p_i f(z) + (1 - p_i) \tilde{f}_i(y \mid s), \quad 1 \leq i \leq m - 1,$$

$$\tilde{f}_m(y \mid s), \quad i = m.$$

- Idea: take  $\tilde{f}_i$  to produce large values of  $Z_i$ .

#### The Conditional Mixture Algorithm

- We can take  $\tilde{f}_i$ 's to be the conditional distributions of the form

$$\tilde{f}_i(z \mid s) = \frac{f(z)I\{z > a(n - s)\}}{\mathbb{P}\{Z > a(n - s)\}}, \quad i \leq m - 1,$$

$$\tilde{f}_m(z \mid s) = \frac{f(z)I\{z > n - s\}}{\mathbb{P}\{Z > n - s\}},$$

where  $a \in (0, 1)$ .

- Note:  $\tilde{f}_i$  is the conditional distribution of  $Z$  given that  $Z$  is large.

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<sup>2</sup>This algorithm is due to [4]

## 2.2 Performance of the Conditional Mixture Algorithm

### The Performance of the Conditional Mixture Algorithm

- The normalized second moment is

$$\begin{aligned} \frac{E[\widehat{p}_n]}{p_n^2} &= \frac{1}{p_n^2} \int \frac{dF}{d\widetilde{F}}(z_1, \dots, z_m) I\{s_m > n\} F(dz_1, \dots, dz_m) \\ &\sim \frac{1}{m^2 \mathbb{P}\{Z_1 > n\}^2} \int \frac{dF}{d\widetilde{F}}(z_1, \dots, z_m) I\{s_m > n\} F(dz_1, \dots, dz_m) \\ &\sim \frac{1}{m^2} \int \frac{1}{\mathbb{P}\{Z_1 > n\}} \frac{dF}{d\widetilde{F}}(nz_1, \dots, nz_m) I\{s_m > 1\} \frac{F(ndz_1, \dots, ndz_m)}{\mathbb{P}\{Z_1 > n\}}. \end{aligned}$$

### The Performance of the Conditional Mixture Algorithm

- Weak convergence:

$$\frac{F(n \cdot)}{\mathbb{P}\{Z_1 > n\}} \rightarrow \sum_{k=1}^m \int I\{ze_k \in \cdot\} \mu_\alpha(dz)$$

- The normalized likelihood ratio is bounded:

$$\sup_n \frac{1}{\mathbb{P}\{Z_1 > n\}} \frac{dF}{d\widetilde{F}}(nz_1, \dots, nz_m) I\{s_m > 1\} < \infty.$$

### The Performance of the Conditional Mixture Algorithm

- The above enables us to show that the normalized second moment converges:

$$\lim_n \frac{E[\widehat{p}_n]}{p_n^2} = \frac{1}{m^2} \left( \sum_{i=1}^{m-1} \frac{a^{-\alpha}}{1-p_i} \prod_{j=1}^{i-1} \frac{1}{p_j} + \prod_{j=1}^{m-1} \frac{1}{p_j} \right).$$

- It is minimized at

$$p_i = \frac{(m-i-1)a^{-\alpha/2} + 1}{(m-i)a^{-\alpha/2} + 1},$$

with minimum

$$\frac{1}{m^2} \left( (m-1)a^{-\alpha/2} + 1 \right)^2.$$

### The Performance of the Conditional Mixture Algorithm

- The conditional mixture algorithm has (almost) vanishing relative error.
- Heavy-tailed heuristics indicate how to design the algorithm.
- Heavy-tailed asymptotics needed to prove efficiency of the algorithm.

### 3 Markov Chain Monte Carlo in Rare-Event Simulation

#### MCMC for Rare Events<sup>3</sup>

- Let  $X$  be random variable with density  $f$  and compute  $p = \mathbb{P}\{X \in A\}$ , where  $A$  is a rare event.
- The zero-variance sampling distribution is

$$F_A(\cdot) = \mathbb{P}\{X \in \cdot \mid X \in A\}, \quad \frac{dF_A}{dx}(x) = \frac{f(x)I\{x \in A\}}{p}.$$

- It is possible to sample from  $F_A$  by constructing a Markov chain with stationary distribution  $F_A$  (e.g. Gibbs sampler or Metropolis-Hastings).
- Idea: sample from  $F_A$  and extract the normalizing constant  $p$ .

#### MCMC for Rare Events

- Let  $X_0, \dots, X_{T-1}$  be a sample of a Markov chain with stationary distribution  $F_A$ .
- Consider a non-negative function  $v(x)$  with  $\int_A v(x)dx = 1$ . The sample mean

$$\hat{q}_T = \frac{1}{T} \sum_{t=0}^{T-1} \frac{v(X_t)I\{X_t \in A\}}{f(X_t)},$$

is an estimator of

$$E_{F_A} \left[ \frac{v(X)I\{X \in A\}}{f(X)} \right] = \int_A \frac{v(x)}{f(x)} \frac{f(x)}{p} dx = \frac{1}{p} \int_A v(x) dx = \frac{1}{p}.$$

- Take  $\hat{q}_T$  as the estimator of  $1/p$ .

#### MCMC for Rare Events

- The rare-event properties of  $\hat{q}_T$  are determined by the choice of  $v$ .
- The large sample properties of  $\hat{q}_T$  are determined by the ergodic properties of the Markov chain.

#### The Normalized Variance

- The rare-event efficiency ( $p$  small) is determined by the normalized variance:

$$\begin{aligned} & p^2 \text{Var}_{F_A} \left( \frac{v(X)}{f(X)} I\{X \in A\} \right) \\ &= p^2 \left( E_{F_A} \left[ \left( \frac{v(X)}{f(X)} I\{X \in A\} \right)^2 \right] - \frac{1}{p^2} \right) \\ &= p^2 \left( \int \frac{v^2(x)}{f^2(x)} \frac{f(x)}{p} dx - \frac{1}{p^2} \right) \\ &= p \int_A \frac{v^2(x)}{f(x)} dx - 1. \end{aligned}$$

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<sup>3</sup>This part is based on [5]

### The Optimal Choice of $v$

- It is optimal to take  $v$  as  $f(x)/p$ . Indeed, then

$$p^2 \text{Var}_{FA} \left( \frac{v(X)}{f(X)} I\{X \in A\} \right) = p \int_A \frac{v^2(x)}{f(x)} dx - 1 = 0.$$

- Insight: take  $v$  as an approximation of the conditional density given the event.

### 3.1 Efficient Sampling for a Heavy-Tailed Random Walk

#### Efficient Sampling for a Heavy-Tailed Random Walk

- Let  $\{Z_k\}$  be iid with density  $f$ .
- Put  $S_n = Z_1 + \dots + Z_n$ .
- Compute  $\mathbb{P}\{S_n > a_n\}$ .
- Heavy-tailed assumption:

$$\lim_{n \rightarrow \infty} \frac{\mathbb{P}\{S_n > a_n\}}{\mathbb{P}\{M_n > a_n\}} = 1.$$

(includes subexponential distributions, see [3]).

#### Gibbs Sampling of a Heavy-Tailed Random Walk

1. Start from a state with  $Z_1 > a_n$ .
2. Update the steps in a random order according to  $j_1, \dots, j_n$  (uniformly without replacement).
3. Update  $Z_{j_k}$  by sampling from

$$\mathbb{P}\{Z \in \cdot \mid Z + \sum_{i \neq j_k} Z_i > n\}.$$

- The vector  $(Z_{t,1}, \dots, Z_{t,n})'$  forms a uniformly ergodic Markov chain with stationary distribution

$$FA(\cdot) = \mathbb{P}\{(Z_1, \dots, Z_n)' \in \cdot \mid Z_1 + \dots + Z_n > a_n\}.$$

#### Rare-Event Efficiency The Choice of $v$

- Take  $v$  to be the density of  $\mathbb{P}\{(Z_1, \dots, Z_n)' \in \cdot \mid M_n > a_n\}$ .
- Then

$$\frac{v(z_1, \dots, z_n)}{f(z_1, \dots, z_n)} = \frac{1}{\mathbb{P}\{M_n > a_n\}} I\{\sum_{i=1}^n z_i > a_n\}.$$

- Rare event efficiency:

$$\begin{aligned} p_n^2 \text{Var}_{FA} \left( \frac{v(Z_1, \dots, Z_n)}{f(Z_1, \dots, Z_n)} I\{Z_1 + \dots + Z_n > a_n\} \right) \\ &= \frac{\mathbb{P}\{S_n > a_n\}^2}{\mathbb{P}\{M_n > a_n\}^2} \mathbb{P}\{M_n > a_n \mid S_n > a_n\} \mathbb{P}\{M_n \leq a_n \mid S_n > a_n\} \\ &= \frac{\mathbb{P}\{S_n > a_n\}}{\mathbb{P}\{M_n > a_n\}} \left( 1 - \frac{\mathbb{P}\{M_n > a_n\}}{\mathbb{P}\{S_n > a_n\}} \right) \rightarrow 0. \end{aligned}$$

$b = 25, T = 10^9, \alpha = 2, n = 5, a = 5, p_{\max} = 0.737e-2$			
	MCMC	IS	MC
Avg. est.	1.050e-2	1.048e-2	1.053e-2
Std. dev.	3e-5	9e-5	27e-5
Avg. time per batch(s)	12.8	12.7	1.4
$b = 25, T = 10^9, \alpha = 2, n = 5, a = 20, p_{\max} = 4.901e-4$			
	MCMC	IS	MC
Avg. est.	5.340e-4	5.343e-4	5.380e-4
Std. dev.	6e-7	13e-7	770e-7
Avg. time per batch(s)	14.4	13.9	1.5
$b = 20, T = 10^9, \alpha = 2, n = 5, a = 10^3, p_{\max} = 1.9992e-7$			
	MCMC	IS	
Avg. est.	2.0024e-7	2.0027e-7	
Std. dev.	3e-11	20e-11	
Avg. time per batch(s)	15.9	15.9	
$b = 20, T = 10^9, \alpha = 2, n = 5, a = 10^4, p_{\max} = 1.99992e-9$			
	MCMC	IS	
Avg. est.	2.00025e-9	2.00091e-9	
Std. dev.	7e-14	215e-14	
Avg. time per batch(s)	15.9	15.9	

$b = 25, T = 10^9, \alpha = 2, n = 20, a = 20, p_{\max} = 1.2437e-4$			
	MCMC	IS	MC
Avg. est.	1.375e-4	1.374e-4	1.444e-4
Std. dev.	2e-7	3e-7	492e-7
Avg. time per batch(s)	52.8	50.0	2.0
$b = 25, T = 10^9, \alpha = 2, n = 20, a = 200, p_{\max} = 1.2494e-6$			
	MCMC	IS	MC
Avg. est.	1.2614e-6	1.2615e-6	1.2000e-6
Std. dev.	4e-10	12e-10	33,166e-10
Avg. time per batch(s)	49.4	48.4	1.9
$b = 20, T = 10^9, \alpha = 2, n = 20, a = 10^3, p_{\max} = 4.9995e-8$			
	MCMC	IS	
Avg. est.	5.0091e-8	5.0079e-8	
Std. dev.	7e-12	66e-12	
Avg. time per batch(s)	53.0	50.6	
$b = 20, T = 10^9, \alpha = 2, n = 20, a = 10^4, p_{\max} = 5.0000e-10$			
	MCMC	IS	
Avg. est.	5.0010e-10	5.0006e-10	
Std. dev.	2e-14	71e-14	
Avg. time per batch(s)	48.0	47.1	

**Numerical illustration**  $n = 5, \alpha = 2, a_n = an.$

**Numerical Illustration**  $n = 20, \alpha = 2, a_n = an.$

### 3.2 Efficient Sampling of Heavy-Tailed Random Sums

#### Efficient Sampling for a Heavy-Tailed Random Sum

- Let  $\{Z_k\}$  be iid with density  $f.$
- Put  $S_n = Z_1 + \dots + Z_n$  and let  $\{N_n\}$  be independent of  $\{Z_k\}$



- Compute  $\mathbb{P}\{S_{N_n} > a_n\}$ .
- Heavy-tailed assumption:<sup>4</sup>

$$\lim_{n \rightarrow \infty} \frac{\mathbb{P}\{S_{N_n} > a_n\}}{\mathbb{P}\{M_{N_n} > a_n\}} = 1.$$

### Constructing the Gibbs Sampler

- Suppose we are in state  $(N_t, Z_{t,1}, \dots, Z_{t,N_t}) = (k_t, z_{t,1}, \dots, z_{t,k_t})$ .
- Let  $k_t^* = \min\{j : z_{t,1} + \dots + z_{t,j} > a_n\}$ .
- Update the number of steps  $N_{t+1}$  from the distribution

$$p(k_{t+1} | k_t^*) = \mathbb{P}\{N = k_{t+1} | N \geq k_t^*\}$$

- If  $k_{t+1} > k_t$ , sample  $Z_{t+1,k_t+1}, \dots, Z_{t+1,k_{t+1}}$  independently from  $F_Z$ .
- Proceed by updating all steps as before.
- Permute the steps.

### Properties of the algorithm

- The Markov chain  $\{(N_t, Z_{t,1}, \dots, Z_{t,N_t})\}$  is uniformly ergodic with stationary distribution

$$F_A(\cdot) = \mathbb{P}\{(Z_1, \dots, Z_n)' \in \cdot \mid Z_1 + \dots + Z_n > a_n\}.$$

- With  $v$  as the density of  $\mathbb{P}\{(N, Z_1, \dots, Z_N) \in \cdot \mid M_N > a_n\}$  we have

$$\begin{aligned} \frac{v(k, z_1, \dots, z_k)}{f(k, z_1, \dots, z_k)} &= \frac{1}{\mathbb{P}\{M_N > a_n\}} I\{\max z_1, \dots, z_k > a_n\} \\ &= \frac{1}{1 - g_N(F_Z(a_n))} I\{\max z_1, \dots, z_k > a_n\}. \end{aligned}$$

- The algorithm has vanishing normalized variance.

**Numerical illustration Geometric( $\rho$ ) number of steps:**  $\rho = 0.05, \alpha = 1, a_n = a/\rho$ .

**Numerical Illustration Geometric( $\rho$ ) number of steps:**  $\rho = 0.05, \alpha = 1, a_n = a/\rho$ .

### References

- [1] S. Asmussen and P.W. Glynn Stochastic Simulation: Algorithms and Analysis. Springer-Verlag, New York, 2007.
- [2] S. Resnick Heavy-Tail Phenomena: Probabilistic and Statistical Modeling. Springer, New York, 2006.
- [3] D. Denisov, A.B. Dieker, and V. Shneer Large deviations for random walks under subexponentiality: the big-jump domain. Ann. Probab. 36(5), 1946-1991, 2008.
- [4] P. Dupuis, K. Leder, and H. Wang Importance sampling for sums of random variables with regularly varying tails. ACM Trans. Model. Comput. Simul. 17(3), 2007.

<sup>4</sup>See e.g. [6] for examples.

$b = 25, T = 10^5, \alpha = 1, \rho = 0.2, a = 10^2, p_{\max} = 0.990e-2$			
	MCMC	IS	MC
Avg. est.	1.149e-2	1.087e-2	1.089e-2
Std. dev.	4e-5	6e-5	35e-5
Avg. time per batch(s)	25.0	11.0	1.2
$b = 25, T = 10^5, \alpha = 1, \rho = 0.2, a = 10^3, p_{\max} = 0.999e-3$			
	MCMC	IS	MC
Avg. est.	1.019e-3	1.012e-3	1.037e-3
Std. dev.	1e-6	3e-6	76e-6
Avg. time per batch(s)	25.8	11.1	1.2
$b = 20, T = 10^6, \alpha = 1, \rho = 0.2, a = 5 \cdot 10^7, p_{\max} = 2.000000e-8$			
	MCMC	IS	
Avg. est.	2.000003e-8	1.999325e-8	
Std. dev.	6e-14	1114e-14	
Avg. time per batch(s)	385.3	139.9	
$b = 20, T = 10^6, \alpha = 1, \rho = 0.2, a = 5 \cdot 10^9, p_{\max} = 2.0000e-10$			
	MCMC	IS	
Avg. est.	2.0000e-10	1.9998e-10	
Std. dev.	0	13e-14	
Avg. time per batch(s)	358.7	130.9	

$b = 25, T = 10^5, \alpha = 1, \rho = 0.05, a = 10^3, p_{\max} = 0.999e-3$			
	MCMC	IS	MC
Avg. est.	1.027e-3	1.017e-3	1.045e-3
Std. dev.	1e-6	4e-6	105e-6
Avg. time per batch(s)	61.5	44.8	1.3
$b = 25, T = 10^5, \alpha = 1, \rho = 0.05, a = 5 \cdot 10^5, p_{\max} = 1.9999e-6$			
	MCMC	IS	MC
Avg. est.	2.0002e-6	2.0005e-6	3.2000e-6
Std. dev.	1e-10	53e-10	55,678e-10
Avg. time per batch(s)	60.7	45.0	1.3

[5] T. Gudmundsson and H. Hult Markov chain Monte Carlo for computing rare-event probabilities for a heavy-tailed random walk. <http://arxiv.org/pdf/1211.2207v1.pdf>

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